

# Bi-Conjugate Gradient Algorithm for Solution of Integral Equations Arising in Electromagnetic Scattering Problems

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# Naval Air Warfare Center Weapons Division

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## FOREWORD

This report details the implementation of applying the bi-conjugant gradient algorithm to solve a volume integral equation. The integral equation is a solution to Maxwell's equations and involves scattering of a plane wave from a metallic nano-cylinder. The solution algorithm is outlined, a numerical example is given, and a FORTRAN code listing is provided. This work was done during a 2-month period at the Naval Air Warfare Center Weapons Division (NAWCWD), China Lake, California, during June and July 2004. The ONR-ASEE Fellowship Program funded it.

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## INTRODUCTION

The work described in this report was performed in collaboration with Professor Surendra Singh, who came to the Naval Air Warfare Center Weapons Division (NAWCWD), China Lake, California, during June and July 2004. Professor Singh, a faculty member in Electrical Engineering Department at The University of Tulsa, Tulsa, Oklahoma, was visiting as an Office of Naval Research-ASEE Summer Faculty Fellow. He worked in the Optics, RF and Material Physics Branch, Sensor and Signal Sciences Division, Research Department. This work involved learning and implementing the bi-conjugate gradient algorithm in order to numerically solve linear equation systems resulting from integral equations arising in electromagnetic scattering problems.

The bi-conjugate gradient (bi-CG) algorithm (Reference 1) is a well-known and established iterative technique to solve a system of linear equations of the form  $\mathbf{A} \mathbf{x} = \mathbf{b}$ . In this equation the square matrix  $\mathbf{A}$  and the source or excitation vector  $\mathbf{b}$  are both complex. The algorithm converges most rapidly if the matrix  $\mathbf{A}$  is square, symmetric, and positive-definite. This does not imply that the algorithm is not suitable if the matrix is not symmetric. We have applied the algorithm to non-symmetric matrices with good results. Details of the bi-CG algorithm are provided in the next section. A variant of this algorithm, the bi-CG stabilized (bi-CGSTAB) algorithm (Reference 2), was also implemented. But in the numerical experimentation, we found that the stabilized version did not provide any advantage over the earlier version. A very good source for a basic understanding of the conjugate gradient method is provided in Reference 3. In the Bi-Conjugate Gradient Method section, we provide the solution of an integral equation known as the Lippman-Schwinger equation. The 2-dimensional physical problem involves a plane wave incident on a cylinder. Note that the original computer code to obtain the numerical solution of the integral equation was developed by Klaus Halterman and J. Merle Elson at China Lake. This code utilized matrix inversion to solve the system of equations. The code was modified to accommodate the bi-CG algorithm. A listing of the FORTRAN code of the bi-CG algorithm and associated function subprograms is provided in the Appendix.

**BI-CONJUGATE GRADIENT METHOD (bi-CG)**

The bi-CG method is an iterative technique used to solve a system of linear equations of the form  $\mathbf{A} \mathbf{x} = \mathbf{b}$ . The method is most suitable in situations where matrix  $\mathbf{A}$  is so large that storing the matrix may pose a significant storage problem. This means that the method can be implemented without storing the entire  $\mathbf{A}$  matrix. This is accomplished by just calculating a single row or column of the matrix at a time. This provides tremendous savings in storage requirements and allows the user to solve a very large system of equations. Note that the savings in storage requirements comes at the cost of increased computation time because each iteration of the algorithm requires computing the elements of matrix  $\mathbf{A}$  twice. In order to increase the computational efficiency in evaluating the matrix elements, we found that the matrix elements involve the computation of two types of Hankel functions. The two Hankel function values are computed once and stored in two arrays. The arrays are then utilized in computing the matrix elements. This technique provided considerable saving in computation time. It is observed that the bi-CG algorithm converges a little slower when the matrix  $\mathbf{A}$  is dense, and that the convergence is faster when the matrix  $\mathbf{A}$  is sparse. The pseudo-code for the bi-CG method is given below:

Initialize the following variables:  $\text{conv} = 10^{-4}$ ,  $\text{rerr} = 100$ ,  $\beta = 0$ , and the following vectors:  $\mathbf{r} = \mathbf{b}$ ,  $\mathbf{m} = \mathbf{0}$ ,  $\mathbf{x} = \mathbf{0}$ ,  $\mathbf{p} = \mathbf{0}$ ,  $\mathbf{pn} = \mathbf{0}$ ,  $t = \mathbf{0}$ .

```

 $\mathbf{m} = \mathbf{r}^*$  (* denotes complex conjugate)
 $\text{val1} = \mathbf{b}^* \mathbf{b}$ 

do  $i = 1, \text{nitm}$ 
    if  $\text{rerr} > \text{conv}$ 
         $\text{val2} = \mathbf{r}^* \mathbf{r}$ 
         $\text{rerr} = \text{abs}(\text{val2} / \text{val1})$ 
         $\text{val3} = \mathbf{m}^* \mathbf{r}$ 
        if ( $i \neq 1$ )  $\beta = \text{val3} / \text{val5}$ 
         $\mathbf{p} = \mathbf{r} + \beta \mathbf{p}$ 
         $\mathbf{pn} = \mathbf{m} + \beta \mathbf{pn}$ 
         $\mathbf{t} = \mathbf{A} \mathbf{p}$  (This operation can be performed by computing one
row of  $\mathbf{A}$  at a time, thus avoiding the storage of matrix  $\mathbf{A}$ .)

         $\text{val4} = \mathbf{pn}^* \mathbf{x}$ 
         $\alpha = \text{val3} / \text{val4}$ 
         $\mathbf{x} = \mathbf{x} + \alpha \mathbf{p}$ 

```



```

         $r = r - \alpha t$ 
         $t = \text{transpose}(A^*) \text{ pn}$  (This operation can be performed by
computing one column of the matrix A at a time, thus avoiding the storage of the matrix
A.)
         $rn = rn - \alpha t$ 
        val5 = val3
    end if
end do

```

#### **SUBROUTINE bcg (b,nunks,nitm,conv,ci,nit,rerr)**

This subroutine iteratively solves a system of linear equations of the form  $[A][ci]=[b]$  using the bi-CG method. The subroutine does not require the storage of the matrix **A**. Rather it uses two function programs, *arow* and *acol* which provide a specific row or column of the matrix, respectively. Note that vector **x** is replaced by **ci**.

#### **Input Variables:**

*nunks* (integer) - Total number of unknowns

*b* (complex vector) - Righthand side or excitation vector of length *nunks*

*nitm* (integer) - Maximum number of iterations for bi-CG (typically set it equal to *nunks*)

*conv* (real) - Convergence factor for bi-CG (typically  $10^{-4}$  to  $10^{-6}$ )

#### **Output Variables:**

*ci* (complex) - Solution vector of length *nunks*

*icount* (integer) - Number of iterations needed for bi-CG method to converge

*rerr* (real) - Residual error. The bi-CG algorithm will stop when the residual error (*rerr*) becomes less than the pre-specified convergence factor (*conv*).

#### **FUNCTION arow(i, vec)**

This function provides the *ith* row of matrix **A** and returns a vector, *vec*, of length *nunks*.



Input:  $i$  (integer) –  $i$ th row number of matrix  $A$

Output:  $vec$  (complex vector) –  $i$ th row of matrix  $A$  vector of length,  $nunkns$ .

### FUNCTION acol (j, vec)

This function provides the  $j$ th column of matrix,  $A$ , and returns a vector,  $vec$ , of length  $nunkns$ .

Input:  $j$  (integer) –  $j$ th row number of matrix  $A$

Output:  $vec$  (complex vector) –  $j$ th column of matrix  $A$  vector of length,  $nunkns$ .

### COMPUTATION OF HANKEL FUNCTIONS $H_0$ AND $H_1$

The tensor Green's function components require the evaluation of Hankel functions  $H_0$  and  $H_1$ . Because the argument of the Hankel functions depends on the relative distance between the source  $(x', y')$  and observation  $(x, y)$  point, it was determined that there are only a limited number of observation and source point combinations that are repeated over and over again in the computation of the Green's function. The Hankel functions are computed for different combinations of source and observation points and stored in two arrays. Then the stored values are utilized in the computation of the tensor Green's function, thereby reducing computation time by avoiding repeated computations.

### APPLICATION OF bi-CG METHOD TO SOLVE AN INTERGAL EQUATION

Here we provide an example of solving a system of linear equations using the bi-CG method. The system of equations is obtained as part of the numerical solution of an integral equation. The first attempt at solving the equations was done by using the traditional approach of matrix inversion. However this required storing the matrix  $A$ . As the dimensionality of the problem increased, thereby increasing the number of unknowns as well as the size of matrix  $A$ , it was clear that the resulting problem could not be solved because of the limited storage capacity of the personal computer: hence the move to bi-CG method.

Consider a scattering system described by a dielectric function,  $\epsilon(r)$ , embedded in an infinite homogeneous background material,  $\epsilon_B$ . When the system is illuminated by an incident field,  $E^0(r)$ , the total electric field,  $E(r)$ , is given by the following integral equation (Reference 4):

$$E(r) = E^0(r) + \int_V dr' G^B(r, r') k_0^2 \Delta\epsilon(r') E(r')$$

where  $G^B(r, r')$  is the Green's tensor associated with the infinite background  $\epsilon_B$ ,  $\Delta\epsilon(r) = \epsilon(r) - \epsilon_B$  and the integration is over the entire scatterer volume,  $V$ . To implement a numerical solution to the above integral equation, we define a grid with  $N$  meshes. Each mesh,  $i$ , is centered at position,  $r_i$ , and has a volume,  $V_i, i = 1, 2, \dots, N$  (for a 2-dimensional systems,  $V_i$  will be replaced by the area of the mesh,  $i$ ). Representing the discretized electric field,  $E_i = E(r_i)$ , the dielectric constant,  $\Delta\epsilon_i = \Delta\epsilon(r_i)$ , and the Green's function,  $G^B(i, j) = G^B(r_i, r_j)$ , the integral equations can be written as a system of linear equations:

The Green's tensor  $G^B(r, r')$  is given by

$$G^B(r, r') = \begin{bmatrix} [G_{xx}^B] & [G_{xy}^B] & [G_{xz}^B] \\ [G_{xy}^B] & [G_{yy}^B] & [G_{yz}^B] \\ [G_{xz}^B] & [G_{yz}^B] & [G_{zz}^B] \end{bmatrix}$$

For a complete description of the Green's tensor components, refer to Reference 4. For a geometry with  $N$  meshes,  $N = nx * ny$ , where  $nx$  and  $ny$  represent the number of divisions of the scatterer in  $x$  and  $y$  directions, respectively, the total number of unknowns is  $3 * N$  to account for the three components of the electric field. Each of the sub-matrices in the tensor Green's function is now of the order of  $N \times N$ . Now we provide a 2-dimensional scattering problem that uses the above formulation to compute the total electric field.

## NUMERICAL EXAMPLE

We provide a numerical example (Reference 5) utilizing the formulation presented in this Section. We consider a metal wire of radius =  $10 \text{ nm}$  with a dielectric constant,  $\epsilon = (-1.07, 0.29)$ . The wavelength of the incident field is  $\lambda = 388 \text{ nm}$ . The geometry is discretized in a grid of  $250 \times 250$  ( $nx = ny = 250$ ), resulting in a total of  $125,000$  ( $2 * nx * ny$ ) unknowns. The number of unknowns in this example is  $2 * nx * ny$ , as only two components of the electric field are needed due to the polarization of the incident field. (Note that if we needed to store the matrix  $\mathbf{A}$  of dimension  $(125,000 \times 125,000)$ , the storage space needed would be very large. This provides the motivation for using the bi-



CG iterative method, in which case we do not store this matrix). Figure 1 shows the electric field amplitude relative to the incident field around the wire as it is illuminated with a plane wave.

### Contour Plot of Relative Field Amplitude

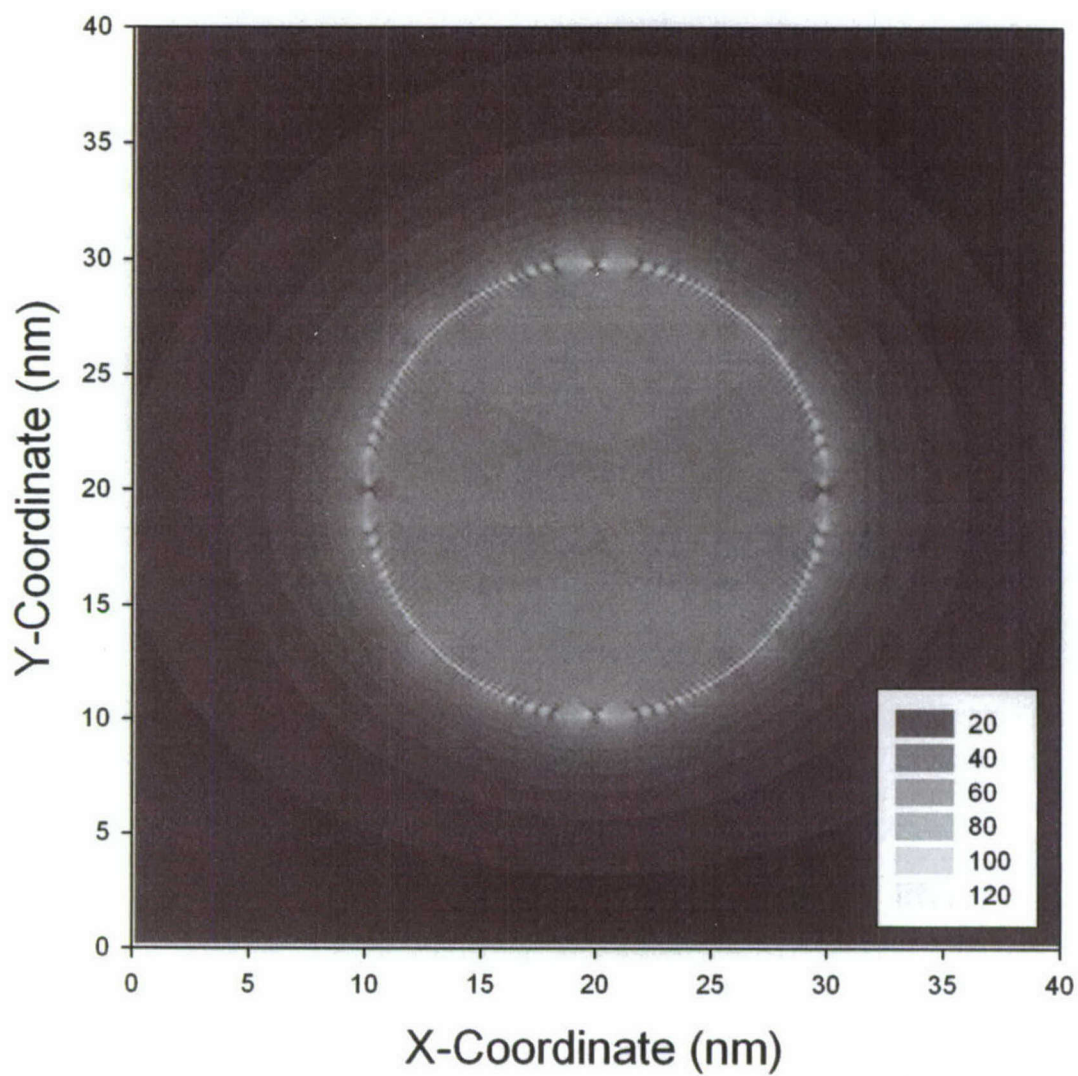


FIGURE 1. Contour Plot of Relative Field Amplitude.



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# **Appendix** **FORTTRAN CODE OF THE BI-CG ALGORITHM AND ASSOCIATED** **FUNCTION SUBPROGRAMS**

Here is the listing of the computer code with subroutine *bcg* and function subprograms *arow* and *acol*:

```

! -----
!
!   bi-conjugant gradient (works)
!   solution to integral equation for e-field
!
! -----
!
!   program bcg_method
!   implicit none
!   real,parameter :: system_size=40 !total grid size in nm
!   real, parameter :: lambda0=338 !incident wavelength in nm
!   real, parameter :: L=system_size/lambda0 !total system size in
wavelengths
!   integer, parameter :: nx=60, ny=nx, nxy = nx*ny !total number of grid
points
!   integer, parameter :: maxi=3*nxy
!   complex, parameter :: Im=(0., 1.),zero=(0., 0.),one=(1., 0.)
!   complex, dimension(1:3*nxy) :: E0,etotal
!   integer, dimension(1:3*nxy) :: ipiv
!   complex, dimension(1:2*nxy) :: H0,H1
!   integer, dimension(1:nxy) :: ix,jy
!   complex, dimension(1:nxy) :: eps
!   real:: arg, r_eff,theta,psi,krho,phi,pi,kx,ky,x,y,wc,wcda,dx,dy
!   real :: xd,yd,yp,yp
!   complex :: mfactor,Mx,My,Mz,kz,eb,z1,mxx,myy,mzz
!   integer :: i,j,k,icount,ij
!   real::conv,rerr,rho1,rho2,rho3,rho
!   INTEGER :: clock_start,clock_end,clock_rate
!   REAL :: elapsed_time,f
!   real :: r1
!   CALL SYSTEM_CLOCK(COUNT_RATE=clock_rate) ! Find the rate
!   CALL SYSTEM_CLOCK(COUNT=clock_start) ! Start timing
!   conv=.0001
!
!   E0=0.
!   eb=one
!   dx=L/nx
!   dy=dx
!   !dy=.025      !grid size per wavelength
!   pi = acos(-1.)
!   wc = 2.*pi

```

```

wcda = wc**2*dx*dy
theta=90.
! phi = 90.
! psi=0.
kx=-cos(theta*pi/180.)
ky=-sin(theta*pi/180.)
kz=0.
krho=sqrt(eb) !krho/k0
r_eff = sqrt(dx*dy/pi)
mfactor=wc*r_eff*H1_fun(wc*r_eff*krho)/krho+2*Im/(pi*krho**2)
Mx = Im*pi/4*(2.-krho**2)*mfactor
My=Mx
Mz = Im*pi/2*(1.-kz**2)*mfactor

k = 0
do i = 1, nx
  do j = 1, ny
    k = k + 1 !k goes up to ny*nx
    x=i*dx
    y=j*dy
    ix(k) = i
    jy(k) = j
    arg=kx*x+ky*y
    z1=exp(wc*Im*arg)
    ! E0(k) =
z1*(cos(phi*pi/180.)*sin(theta*pi/180.)*cos(psi*pi/180.)-&
! sin(phi*pi/180.)*sin(psi*pi/180.))
! E0(k+nxy) =
z1*(sin(phi*pi/180.)*sin(theta*pi/180.)*cos(psi*pi/180.)+&
! cos(phi*pi/180.)*sin(psi*pi/180.))
!E0(k+2*nxy) = z1*cos(theta*pi/180.)* cos(psi*pi/180.)
E0(k+2*nxy) = z1
end do
end do

eps = zero
r1=10./lambda0 !radius in wavelengths
! choose sites of delta_epsilon over 2D grid
k = 0
do i = 1,nx
  do j = 1,ny
    k = k + 1
    eps(k)=zero
    ! eps(k) = one*1.5
    rho1=sqrt(real(i-nx/2)**2+real(j-ny/2)**2)
    if (rho1<=int(r1/dx)) eps(k) = (-1.07,0.29) - one
    rho2=sqrt(real(i-nx/2)**2+real(j-ny/2)**2)
    ! if (rho2<=3.) eps(k) = (-100.,0.) - one
    rho3=sqrt(real(i-3*(nx/4))**2+real(j-ny/2)**2)
    ! if (rho3<=3.) eps(k) = (-100.,0.) - one
    ! if(i >= nx/2-5 .and. i <= nx/2+5 .and. &
    ! j >= ny/2-5 .and. j <= ny/2+5) eps(k) = (-100.,0.) - one
    ! write(45,445) ix(k),jy(k),real(eps(k))
end do
end do
!close(45)

i=1
do j=2,nxy

```



```

x = ix(i)*dx;y=jy(i)*dy
  xp = ix(j)*dx;yp = jy(j)*dy
  xd=x-xp;yd=y-yp
  rho = sqrt(xd*xd+yd*yd)
  arg = wc*krho*rho
  ij=(ix(i)-ix(j))**2 + (jy(i)-jy(j))**2
H0(ij)=H0_fun(arg)
H1(ij)=H1_fun(arg)
end do

mx=x-Mx-0.5
my=y-My-0.5
mz=z-Mz
! *****
*****

      call bcg(E0,3*nxy,maxi,conv,etotal,icount,rerr)
      print *, 'number of iterations=', icount-1
      print *, 'the residual error =', rerr
! print *, testx
! *****
*****

      do k = 1, nxy !Ex
        x = ix(k)*dx ; y = jy(k)*dy
        write(97,*) x*lambda0, y*lambda0, abs(etotal(k))
      end do

      do k = nxy+1, 2*nxy !Ey
        x = ix(k-nxy)*dx ; y = jy(k-nxy)*dy
        write(98,*) x*lambda0, y*lambda0, abs(etotal(k))
      end do

      do k = 2*nxy+1, 3*nxy !Ez
        x = ix(k-2*nxy)*dx ; y = jy(k-2*nxy)*dy
        write(99,*) x*lambda0, y*lambda0, abs(etotal(k))
      end do

      do k=1,nxy
        x = ix(k)*dx ; y = jy(k)*dy

arg=etotal(k)*conjg(etotal(k))+etotal(k+nxy)*conjg(etotal(k+nxy))+etotal
(k+2*nxy)*conjg(etotal(k+2*nxy))
        write(96,*) x*lambda0, y*lambda0, arg
      end do
! *****
*****

444 format(F6.3,1X,F6.3,1X,F6.3)

CALL SYSTEM_CLOCK(COUNT=clock_end) ! Stop timing
! Calculate the elapsed time in seconds:
elapsed_time=REAL((clock_end-clock_start))/clock_rate
print *, 'elapsed time =', elapsed_time/60., 'minutes'
contains

subroutine bcg(b,nunkns,nitm,conv,ci,nit,rerr)
implicit none
complex, intent(in), dimension(1:nunkns) :: b

```

```

complex, intent(out), dimension(1:nunkns) :: ci
integer, intent(in) :: nunkns, nitm
integer, intent(out) :: nit
real, intent(in) :: conv
real, intent(out) :: rerr
complex, dimension(1:nunkns) :: x, p, pn, r, rn
complex :: beta, alpha, val1, val2, val3, val4, val5, sumx
integer :: i

x=0;p=0;pn=0;ci=0;icount=0;beta=0
rerr=100
r=b
rn=conjg(r)
val1=dot_product(conjg(b),b)
do nit=1,nitm

  if (rerr>conv) then
    val2=dot_product(conjg(r),r)
    rerr=abs(val2/val1)
    val3=dot_product(conjg(rn),r)
    if (nit/=1) beta=val3/val5
    p=r+beta*p
    pn=rn+beta*pn

    forall (i=1:nunkns) x(i)=dot_product(conjg(arow(i)),p)

    val4=dot_product(conjg(pn),x)
    alpha=val3/val4
    ci=ci+alpha*p
    r=r-alpha*x
    forall (i=1:nunkns) x(i)=dot_product(conjg(acol(i)),pn)

    rn=rn-alpha*x
    val5=val3
    icount=icount+1
    print *, 'iteration#', nit-1
    print *, 'residual error', rerr
  else
    exit
  end if

end do
end subroutine bcg

function arow(i)
implicit none
integer, intent(in) :: i
complex, dimension(1:3*nxy) :: arow
integer :: j, n

n=nxy

do j=1,3*n

  if (i<=n) then

    if (j<=n) then
      if (i==j) then
        arow(j)=1.-mxx*eps(j)
      else
        if (eps(j)==zero) then

```

```

        arow(j)=0.
      else
        arow(j)=-gxx(i,j)*eps(j)
      end if
    end if
  else if (j>n .and. j<=2*n) then
    if (eps(j-n)==zero) then
      arow(j)=0.
    else
      arow(j)=-gxy(i,j-n)*eps(j-n)
    end if
  else if (j>2*n) then
    if (eps(j-2*n)==zero) then
      arow(j)=0.
    else
      arow(j)=-gxz(i,j-2*n)*eps(j-2*n)
    end if
  end if
end if

end if

if (i>n .and. i<=2*n) then

  if (j<=n) then
    if (eps(j)==zero) then
      arow(j)=0.
    else
      arow(j)=-gxy(i-n,j)*eps(j)
    end if
  else if (j>n .and. j<= 2*n) then
    if (i==j) then
      arow(j)=1.-myy*eps(j-n)
    else
      if (eps(j-n)==zero) then
        arow(j)=0.
      else
        arow(j)=-gyy(i-n,j-n)*eps(j-n)
      end if
    end if
  else if (j>2*n) then
    if (eps(j-2*n)==zero) then
      arow(j)=0.
    else
      arow(j)=-gyz(i-n,j-2*n)*eps(j-2*n)
    end if
  end if
end if

end if

if (i>2*n) then

  if (j<=n) then
    if (eps(j)==zero) then
      arow(j)=0.
    else
      arow(j)=-gxz(i-2*n,j)*eps(j)
    end if
  else if (j>n .and. j<= 2*n) then
    if (eps(j-n)==zero) then
      arow(j)=0.
    end if
  end if
end if

```



```

        else
            arow(j)=-gyz(i-2*n,j-n)*eps(j-n)
        end if
    else if (j>2*n) then
        if (i==j) then
            arow(j)=1.-mzz*eps(j-2*n)
        else
            if (eps(j-2*n)==zero) then
                arow(j)=0.
            else
                arow(j)=-gzz(i-2*n,j-2*n)*eps(j-2*n)
            end if
        end if
    end if
end if

```

```

end do
end function arow

```

```

function acol(j)
implicit none
integer, intent(in) :: j
complex, dimension(1:3*nxy) :: acol
integer :: n,i

```

```

n=nxy
do i=1,3*n

```

```

    if (i<=n) then
        if (j<=n) then
            if (i==j) then
                acol(i)=1.-mxx*eps(j)
            else
                if (eps(j)==zero) then
                    acol(i)=0.
                else
                    acol(i)=-gxx(i,j)*eps(j)
                end if
            end if
        else if (j>n .and. j<=2*n) then
            if (eps(j-n)==zero) then
                acol(i)=0.
            else
                acol(i)=-gxy(i,j-n)*eps(j-n)
            end if
        else if (j>2*n) then
            if (eps(j-2*n)==zero) then
                acol(i)=0.
            else
                acol(i)=-gxz(i,j-2*n)*eps(j-2*n)
            end if
        end if
    end if
end if

```

```

    if (i>n .and. i<=2*n) then

```

```

        if (j<=n) then
            if (eps(j)==zero) then

```

```

        acol(i)=0.
        else
        acol(i)=-gxy(i-n,j)*eps(j)
        end if

else if (j>n .and. j<= 2*n) then
    if (i==j) then
        acol(i)=1.-myy*eps(j-n)
    else
        if (eps(j-n)==zero) then
            acol(i)=0.
        else
            acol(i)=-gyy(i-n,j-n)*eps(j-n)
        end if
    end if
else if (j>2*n) then
    if (eps(j-2*n)==zero) then
        acol(i)=0.
    else
        acol(i)=-gyz(i-n,j-2*n)*eps(j-2*n)
    end if
end if
end if

if (i>2*n) then

    if (j<=n) then
        if (eps(j)==zero) then
            acol(i)=0.
        else
            acol(i)=-gxz(i-2*n,j)*eps(j)
        end if

    else if (j>n .and. j<= 2*n) then
        if (eps(j-n)==zero) then
            acol(i)=0.
        else
            acol(i)=-gyz(i-2*n,j-n)*eps(j-n)
        end if
    else if (j>2*n) then
        if (i==j) then
            acol(i)=1.-mzz*eps(j-2*n)
        else
            if (eps(j-2*n)==zero) then
                acol(i)=0.
            else
                acol(i)=-gzz(i-2*n,j-2*n)*eps(j-2*n)
            end if
        end if
    end if
end if

end do
end function acol

complex function gxx(i,j)
    implicit none

```

```

integer,intent(in) :: i,j
integer :: ij,dif1,dif2
real :: rho,x,y,yp,rx,ry,c2theta,xd,yd

x = ix(i)*dx;y = jy(i)*dy
xp = ix(j)*dx;yp = jy(j)*dy
xd=x-xp;yd=y-yp
rho = sqrt(xd*xd+yd*yd)
dif1= ix(i)-ix(j);dif2=jy(i)-jy(j)
ij=dif1*dif1+dif2*dif2
rx = xd/rho
ry = yd/rho
c2theta=rx*rx-ry*ry
gxx=wcda*Im/4.*(1.-
krho*rx*krho*rx)*H0(ij)+krho*c2theta*H1(ij)/(wc*rho))

end function gxx

complex function gxy(i,j)
implicit none
integer,intent(in) :: i,j
integer :: ij,dif1,dif2
real :: rho,x,y,yp,rx,ry,s2theta,xd,yd
if(i /= j) then
x = ix(i)*dx;y=jy(i)*dy
xp = ix(j)*dx;yp=jy(j)*dy
xd=x-xp;yd=y-yp
rho = sqrt(xd*xd+yd*yd)
arg = wc*krho*rho
dif1= ix(i)-ix(j);dif2=jy(i)-jy(j)
ij=dif1*dif1+dif2*dif2
rx = xd/rho
ry = yd/rho
s2theta=2*ry*rx
gxy=wcda*Im/8.*krho*krho*s2theta*(2./arg*H1(ij)-H0(ij))
else
gxy = zero
end if
end function gxy

complex function gxz(i,j)
implicit none

integer,intent(in) :: i,j
integer :: ij,dif1,dif2
real :: rho,x,y,yp,rx,xd,yd
if(i /= j) then
x=ix(i)*dx ; y=jy(i)*dy
xp=ix(j)*dx ; yp=jy(j)*dy
xd=x-xp;yd=y-yp
rho = sqrt(xd*xd+yd*yd)
arg = wc*krho*rho
dif1= ix(i)-ix(j);dif2=jy(i)-jy(j)
ij=dif1*dif1+dif2*dif2
rx = xd/rho
gxz=wcda*1./4.*krho*kz*rx*H1(ij)
else
gxz = zero
end if
end function gxz

```



```

complex function gyy(i,j)
  implicit none
  integer,intent(in) :: i,j
  integer :: ij,dif1,dif2
  real :: rho,x,y,yp,rx,ry,c2theta,xd,yd
  x = ix(i)*dx ; y = jy(i)*dy
  xp = ix(j)*dx ; yp = jy(j)*dy
  xd = x - xp ; yd = y - yp
  rho = sqrt(xd*xd + yd*yd)
  arg = wc*krho*rho
  dif1 = ix(i) - ix(j) ; dif2 = jy(i) - jy(j)
  ij = dif1*dif1 + dif2*dif2
  rx = xd/rho
  ry = yd/rho
  c2theta = rx*rx - ry*ry
  gyy = wcda*Im/4.*((1.-krho*ry*krho*ry)*H0(ij) -
krho*c2theta*H1(ij)/(wc*rho))

```

```

end function gyy

```

```

complex function gyz(i,j)
  implicit none
  integer,intent(in) :: i,j
  integer :: ij,dif1,dif2
  real :: rho,x,y,yp,ry,xd,yd
  if(i /= j) then
    x = ix(i)*dx ; y = jy(i)*dy
    xp = ix(j)*dx ; yp = jy(j)*dy
    xd = x - xp ; yd = y - yp
    rho = sqrt(xd*xd + yd*yd)
    arg = wc*krho*rho
    dif1 = ix(i) - ix(j) ; dif2 = jy(i) - jy(j)
    ij = dif1*dif1 + dif2*dif2
    ry = yd/rho
    gyz = wcda*1./4.*krho*kz*ry*H1(ij)
  else
    gyz = zero
  end if
end function gyz

```

```

complex function gzz(i,j)
  implicit none
  integer,intent(in) :: i,j
  integer :: ij,dif1,dif2
  real :: rho,x,y,yp,xd,yd
  x = ix(i)*dx ; y = jy(i)*dy
  xp = ix(j)*dx ; yp = jy(j)*dy
  xd = x - xp ; yd = y - yp
  rho = sqrt(xd*xd + yd*yd)
  arg = wc*krho*rho
  dif1 = ix(i) - ix(j) ; dif2 = jy(i) - jy(j)
  ij = dif1*dif1 + dif2*dif2
  gzz = wcda*Im/4.*((1.-kz*kz)*H0(ij)

```

```

end function gzz

```

```

FUNCTION H0_fun(z) RESULT (H0out)
IMPLICIT NONE
real, INTENT(IN) :: z
real :: x,J_0,Y_0,f_0,theta,x2,x3,x4,x5,x6,x8,x10,x12,sqz

```

```

complex::H0out

if (z <= 3.) then
x=z/3.
x2=x*x;x4=x2*x2;x6=x4*x2;x8=x6*x2;x10=x8*x2;x12=x10*x2
J_0=1.-2.2499997*x2+1.2656208*x4-.3163866*x6+.0444479*x8-
.0039444*x10+.0002100*x12
Y_0=(2/pi)*log(z/2.)*J_0+.36746691+.60559366*x2-
.74350384*x4+.25300117*x6-.04261214*x8+.00427916*x10-.00024846*x12
else
x=3./z
x2=x*x;x3=x2*x;x4=x3*x;x5=x4*x;x6=x5*x;sqz=sqrt(z)
f_0=.79788456-.00000077*x-.00552740*x2-.00009512*x3+.00137237*x4 &
-.00072805*x5 + .00014476*x6

theta=z-.78539816-.04166397*x-.00003954*x2+.00262573*x3-.00054125*x4 &
& -.00029333*x5+.00013558*x6

J_0=f_0*cos(theta)/sqz
Y_0=f_0*sin(theta)/sqz
end if

H0out=cplx(J_0,Y_0)
END FUNCTION H0_fun

FUNCTION H1_fun(z) RESULT (H1out)
IMPLICIT NONE
real, INTENT(IN) :: z
real:: x,J_1,Y_1,f_1,theta1,x2,x3,x4,x5,x6,x8,x10,x12,sqz
complex::H1out

if (z <= 3.) then
x=z/3.
x2=x*x;x4=x2*x2;x6=x4*x2;x8=x6*x2;x10=x8*x2;x12=x10*x2
J_1=.5*z-.56249985*z*x2+.21093573*z*x4-.03954289*z*x6+.00443319*z*x8-
.00031761*z*x10+.00001109*z*x12

Y_1=(2/pi)*log(z/2.)*J_1-.6366198/z+.2212091*x2/z+.2.1682709*x4/z-
1.3164827*x6/z+.3123951*x8/z-.0400976*x10/z+.0027873*x12/z
else
x=3./z
x2=x*x;x3=x2*x;x4=x3*x;x5=x4*x;x6=x5*x;sqz=sqrt(z)
f_1=.79788456+.00000156*x+.01659667*x2+.00017105*x3-
.00249511*x4+.00113653*x5-.00020033*x6

theta1=z-2.35619449+.12499612*x+.00005650*x2-
.00637879*x3+.00074348*x4+.00079824*x5-.00029166*x6

J_1=f_1*cos(theta1)/sqz
Y_1=f_1*sin(theta1)/sqz
end if

H1out=cplx(J_1,Y_1)
END FUNCTION H1_fun

end program bcg_method

```

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